

1. (a) (5, 1, 2) Binary Repetition Code

$$R = \frac{k}{n} = \frac{1}{5} = 0.2$$

$$P_e = 1 - \sum_{k=0}^2 \binom{5}{k} \epsilon^k (1-\epsilon)^{5-k} = 1 - [(1-\epsilon)^5 + 5\epsilon(1-\epsilon)^4 + 10\epsilon^2(1-\epsilon)^3]$$

$$= 10\epsilon^3 - 15\epsilon^4 + 6\epsilon^5$$

(b) (7, 4, 1) Hamming Code

$$R = \frac{k}{n} = \frac{4}{7} \approx 0.5714$$

$$P_e = 1 - \sum_{k=0}^1 \binom{7}{k} \epsilon^k (1-\epsilon)^{7-k} = 1 - [(1-\epsilon)^7 + 7\epsilon(1-\epsilon)^6]$$

$$= 21\epsilon^2 - 70\epsilon^3 + 105\epsilon^4 - 84\epsilon^5 + 35\epsilon^6 - 6\epsilon^7$$

(c) (23, 12, 3) Binary Golay Code

$$R = \frac{12}{23} \approx 0.5217$$

$$P_e = 1 - \sum_{k=0}^3 \binom{23}{k} \epsilon^k (1-\epsilon)^{23-k} = \text{Thank You Mathematica} \dots$$

$$= 1 - (1-\epsilon)^{23} - 23\epsilon(1-\epsilon)^{22} - 253\epsilon^2(1-\epsilon)^{21} - 1771\epsilon^3(1-\epsilon)^{20}$$

$$= 8855\epsilon^4 - 134596\epsilon^5 + 1009470\epsilon^6 - 4903140\epsilon^7$$

$$+ 17160990\epsilon^8 - 45762640\epsilon^9 + 96101544\epsilon^{10}$$

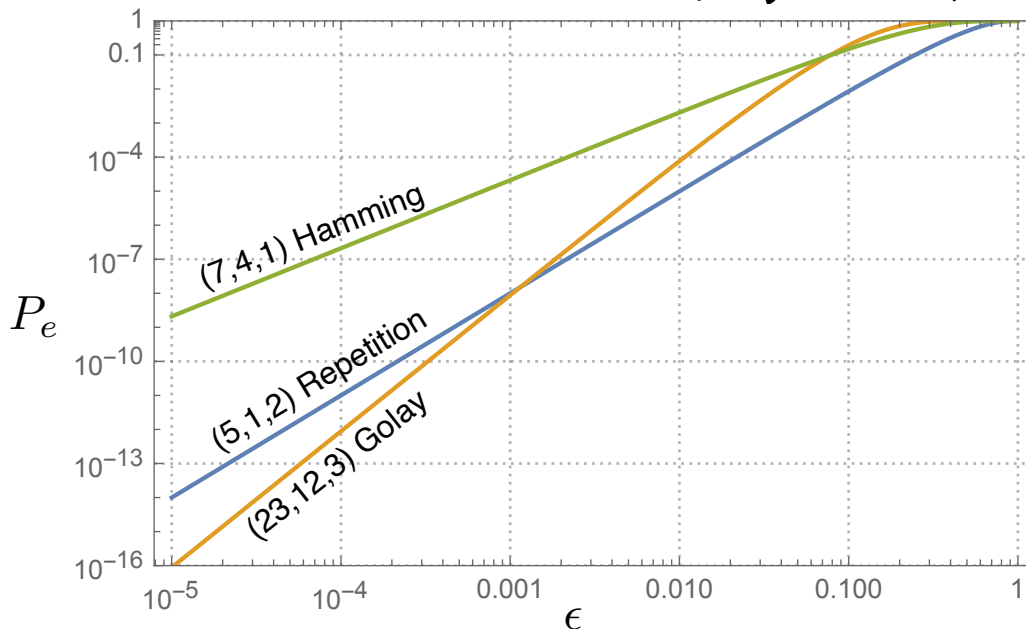
$$- 162249360\epsilon^{11} + 223092870\epsilon^{12} - 251694520\epsilon^{13}$$

$$+ 233716340\epsilon^{14} - 178474296\epsilon^{15} + 111546435\epsilon^{16}$$

$$+ 56530320\epsilon^{17} + 22881320\epsilon^{18} - 7225680\epsilon^{19}$$

$$+ 1716099\epsilon^{20} - 288420\epsilon^{21} + 30590\epsilon^{22} - 1540\epsilon^{23}$$

Plots of these block error probabilities as a function of ϵ for $10^{-5} \leq \epsilon \leq 1$ on a log-log scale appear as follows:



2.(a) A $(2t+1, 1, t)$ binary repetition code can correct any error pattern of t or fewer errors out of $2t+1$ transmitted binary digits. Thus

$$P_e = 1 - \sum_{r=0}^t \binom{2t+1}{r} \epsilon^r (1-\epsilon)^{2t+1-r}$$

(b) A $(k(2t+1), k, t)$ extended binary repetition code is used to transmit k binary information digits using $k(2t+1)$ binary codeword digits. Each codeword is formed by concatenating k $(2t+1, 1, t)$ binary repetition code codewords, one for each of the k binary information digits to be encoded.

Note that this code can correct any pattern of t errors out of $k(2t+1)$ transmitted digits. Note, however, that it cannot correct all patterns of $t+1$ errors, since if $t+1$ errors (for example) occur in the first block of $2t+1$ binary digits, an error will be made. $\therefore S = t$.

In order for a $(k(2t+1), k, t)$ extended repetition code to correct t or less errors, all k sub-blocks of $2t+1$ binary digits (each of these is a $(2t+1, 1, t)$ binary repetition code) must be correctly decoded.

Let $P_e(t, k) \triangleq$ Prob. of block error in $(k(2t+1), k, t)$ extended repetition code.

Then

$$P_e(t, k) = 1 - [1 - P_e(t, 1)]^k$$

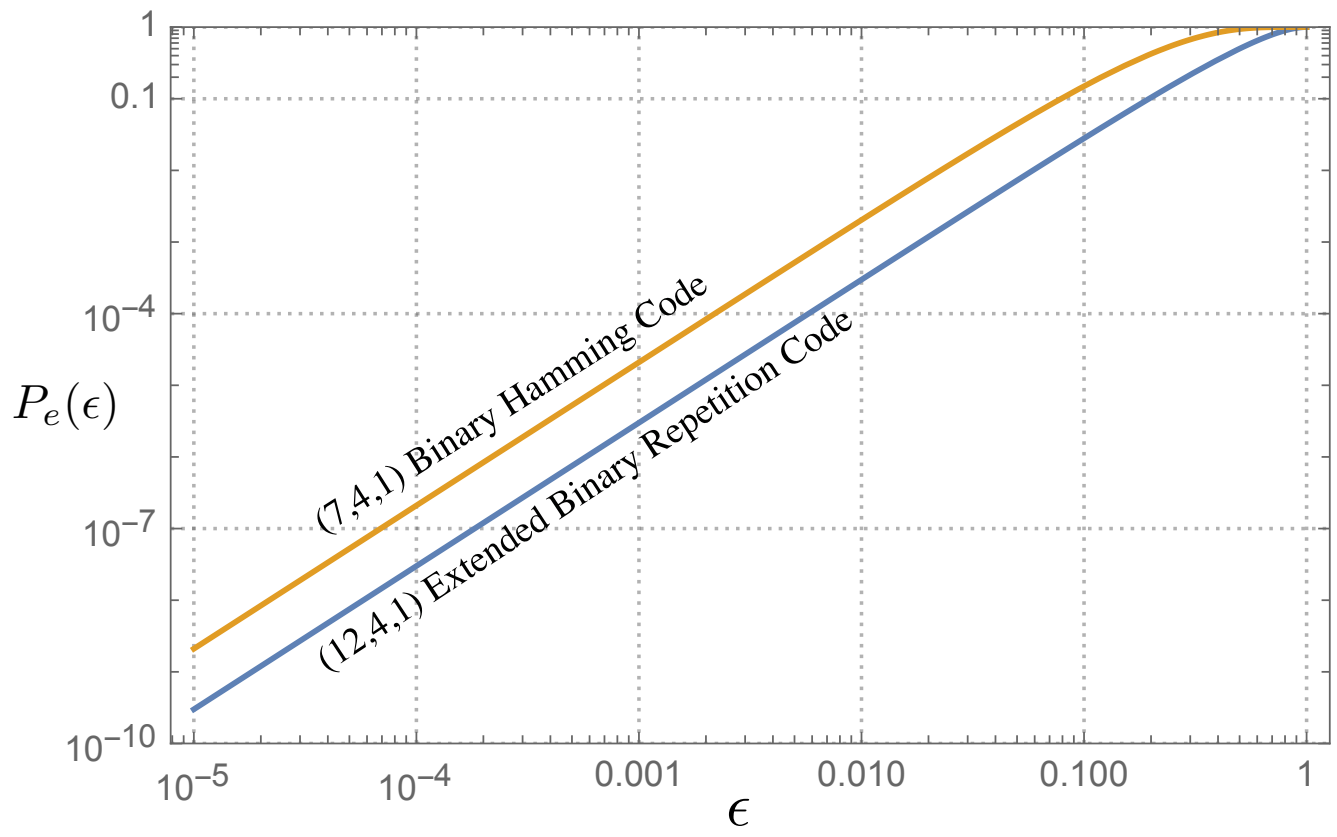
$$= 1 - \left[\sum_{r=0}^t \binom{2t+1}{r} \epsilon^r (1-\epsilon)^{2t+1-r} \right]^k$$

For the purpose of comparison with a $(7, 4, 1)$ Hamming code, if we want our extended repetition code to send 4 information bits, and correct any 1 error, then we should take $k=4$ and $t=1$. So we are comparing a $(12, 4, 1)$ extended binary repetition code to a $(7, 4, 1)$ Hamming code.

$(7, 4, 1)$ Hamming: $R = \frac{k}{n} = \frac{4}{7}$, $P_e = 21\epsilon^2 - 70\epsilon^3 + 105\epsilon^4 - 84\epsilon^5 + 35\epsilon^6 - 6\epsilon^7$

$(12, 4, 1)$ Repetition: $R = \frac{4}{12} = \frac{1}{3}$, $P_e = 1 - (1-\epsilon)^3 - 3(1-\epsilon)^2\epsilon$
 $= 3\epsilon^2 - 2\epsilon^3$

Plots of the error performance of the two codes are given on the following page.



$$\begin{aligned}
3. (a) \quad h(X, Y) - h(X) - h(Y) &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \log \frac{1}{f(x, y)} dx dy \\
&\quad - \int_{-\infty}^{\infty} f(x) \log \frac{1}{f(x)} dx - \int_{-\infty}^{\infty} f(y) \log \frac{1}{f(y)} dy \\
&= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \log \frac{f(x)f(y)}{f(x, y)} dx dy \stackrel{①}{\leq} \log \left(\int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \frac{f(x)f(y)}{f(x, y)} dx dy \right) \\
&= \log \left(\int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x)f(y) dx dy \right) = \log(1) = 0.
\end{aligned}$$

Since $\log(\cdot)$ is strictly convex \cap , equality occurs iff $\frac{f(x)f(y)}{f(x, y)} = \text{const.} = 1 \iff f(x, y) = f(x)f(y)$

$\iff X$ and Y are statistically independent

$\therefore h(X, Y) \leq h(X) + h(Y)$, with equality $\iff X \perp\!\!\!\perp Y$.

(b) Show $h(X) \geq h(X|Y) \iff h(X|Y) - h(X) \leq 0$.

$$\begin{aligned}
h(X|Y) - h(X) &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \log \frac{1}{f(x|y)} dx dy - \int_{-\infty}^{\infty} f(x) \log \frac{1}{f(x)} dx \\
&= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \log \frac{f(x)}{f(x|y)} dx dy = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \log \frac{f(x)f(y)}{f(x, y)} dx dy
\end{aligned}$$

$$\stackrel{①}{\leq} \log \left(\int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x, y) \frac{f(x)f(y)}{f(x, y)} dx dy \right) = \log \left(\int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x)f(y) dx dy \right)$$

$$= \log 1 = 0, \text{ with equality iff } \frac{f(x)f(y)}{f(x, y)} = \text{const} = 1$$

$\iff X$ and Y are statistically independent

$\therefore h(X) \geq h(X|Y)$, with equality iff $X \perp\!\!\!\perp Y$.

$$(c) \quad h(Y, Z|X) - h(Y|X) - h(Z|X) = \iiint_{\mathbb{R}^3} f(x, y, z) \log \frac{f(y|x)f(z|x)}{f(y, z|x)} dx dy dz$$

$$\stackrel{①}{=} \log \left(\iiint_{\mathbb{R}^3} f(x, y, z) \frac{f(y|x)f(z|x)f(x)}{f(x, y, z)} dx dy dz \right)$$

$$= \log \left(\iiint_{\mathbb{R}^3} f(y|x)f(z|x)f(x) dx dy dz \right) = \log \left(\int_{-\infty}^{\infty} f(x) \left[\int_{-\infty}^{\infty} f(y|x) dy \right] \left[\int_{-\infty}^{\infty} f(z|x) dz \right] dx \right)$$

$$= \log \left(\int_{-\infty}^{\infty} f(x) \cdot 1 \cdot 1 dx \right) = \log(1) = 0, \text{ with equality iff}$$

$$\frac{f(y, z|x)}{f(y|x)f(z|x)} = 1 \iff f(y, z|x) = f(y|x)f(z|x) \iff X \text{ and } Y \text{ are stat. indep conditioned on } Z \text{ (} X|Z \perp\!\!\!\perp Y|Z \text{)}.$$

$$4. (a) f_x(x) = \begin{cases} \lambda e^{-\lambda x}, & x \geq 0 \\ 0, & x < 0 \end{cases}$$

$$h(X) = - \int_0^{\infty} \lambda e^{-\lambda x} \log(\lambda e^{-\lambda x}) dx = \dots \stackrel{\text{Integration by parts}}{=} \log e - \log \lambda = \log \frac{e}{\lambda} \\ = 1 - \ln \lambda \text{ (nats)}$$

$$(b) f(x) = \begin{cases} \frac{1}{\pi \sqrt{x(1-x)}}, & 0 < x < 1 \\ 0, & \text{elsewhere} \end{cases} \Rightarrow h(X) = \int_0^1 \frac{1}{\pi \sqrt{x(1-x)}} \log[\pi \sqrt{x(1-x)}] dx$$

This is most easily evaluated using a trigonometric substitution

$$\sqrt{ax^2 + bx + c} = \sqrt{x(1-x)} = \sqrt{x - x^2} \Rightarrow a = -1, b = 1, c = 0$$

From calculus, if you have an expression of the form $\sqrt{-a} \sqrt{\beta^2 - (x-\alpha)^2}$ (which we have with $\alpha = 1/2$ and $\beta = 1/2$), we can make the substitution

$$x - \alpha = \beta \sin \theta \Rightarrow \sqrt{\beta^2 - (x-\alpha)^2} = \sqrt{\beta^2 - \beta^2 \sin^2 \theta} = \beta \cos \theta$$

and $\frac{dx}{d\theta} = \beta \cos \theta$, so we can write $h(X)$ as

$$h(X) = \int_{-\pi/2}^{\pi/2} \frac{1}{\pi \beta \cos \theta} \log[\pi \beta \cos \theta] \beta \cos \theta d\theta = \frac{1}{\pi} \int_{-\pi/2}^{\pi/2} \log[\pi \beta \cos \theta] d\theta \\ = \frac{2}{\pi} \int_0^{\pi/2} \log[\pi \beta \cos \theta] d\theta = \frac{2}{\pi} \int_0^{\pi/2} \log[\pi \beta] d\theta + \frac{2}{\pi} \int_0^{\pi/2} \log[\cos \theta] d\theta \\ = \log \pi \beta + \frac{2}{\pi} \int_0^{\pi/2} \log[\cos \theta] d\theta$$

$$= \log e^{\frac{\pi}{2}} + \frac{2}{\pi} \left(-\frac{\pi}{2} \ln 2\right) \\ = \ln \frac{\pi}{2} - \ln 2 = \ln \frac{\pi}{4} \text{ (nats).}$$

recall: $\beta = 1/2$, and from integral tables

$$\int_0^{\pi/2} \ln[\cos x] dx = -\frac{\pi}{2} \ln 2$$

(c) We can write the triangle density as $f(x) = \begin{cases} \frac{t}{\alpha} x, & x \leq \alpha \\ \frac{2-tx}{2t^{-1}-\alpha}, & \alpha \leq x \leq 2t^{-1} \end{cases}$ (here $\alpha \in [0, 2t^{-1}]$).

Now $\int_0^{\alpha} \frac{t}{\alpha} x \log\left(\frac{t}{\alpha} x\right) dx = \frac{\alpha}{t} \int_0^{\frac{t}{\alpha}} u \log u du$ and $\int_{\alpha}^{2t^{-1}} \left(\frac{2-tx}{2t^{-1}-\alpha}\right) \log\left(\frac{2-tx}{2t^{-1}-\alpha}\right) dx = \dots = \left(\frac{2t^{-1}-\alpha}{t}\right) \int_0^{\frac{t}{2t^{-1}-\alpha}} u \log u du$

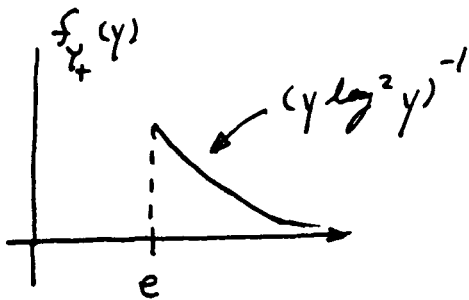
Thus $h(X) = \left[-\frac{\alpha}{t} + \frac{\alpha}{t} - \frac{2t^{-1}}{t}\right] \int_0^{\frac{t}{2t^{-1}-\alpha}} u \log u du = \dots = \frac{1}{2} - \ln 2 \text{ (nats)}$

5. There are many possible solutions. Here's one.

P.5

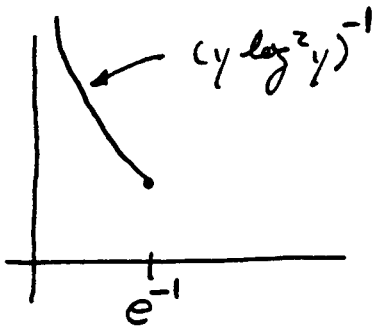
Let X be an r.v. with density $f(x) = \begin{cases} x^{-2}, & x \geq 1 \\ 0, & \text{elsewhere} \end{cases}$
 A simple calculation yields $h(X) = 2$ nats.

(a) Define an r.v. $Y_+ = e^X$. This has a pdf that appears as follows



and a quick calculation shows
 $h(Y_+) = \int_e^\infty \frac{1}{y \log^2 y} \log [y \log^2 y] dy = +\infty$
 i.e. the integral diverges to $+\infty$

(b) Define the r.v. $Y_- = e^{-X}$. The pdf of Y_- appears as follows



and the differential entropy in this case is
 $h(Y_-) = \int_0^{e^{-1}} \frac{1}{y \log^2 y} \log [y \log^2 y] dy = -\infty$
 i.e. the integral diverges to $-\infty$.

6. The easiest way to solve this problem is to use the following Lemma and guess* the appropriate solution

Lemma: Given any two pdf's $f(x)$ and $g(x)$ defined on a subset $A \subseteq \mathbb{R}$ (the real line),

$$-\int_A f(x) \log f(x) dx \leq -\int_A f(x) \log g(x) dx, \quad \text{with equality iff } f(x) = g(x).$$

Proof: $-\int_A f(x) \log f(x) dx + \int_A f(x) \log g(x) dx = \int_A f(x) \log \frac{g(x)}{f(x)} dx \stackrel{J}{\leq} \log \int_A f(x) \frac{g(x)}{f(x)} dx$
 $= \log \int_A g(x) dx = \log 1 = 0$, with equality iff $f(x) = g(x)$, by Jensen's inequality and the strict convexity of $\log(\cdot)$.

* Lagrange Multipliers may be useful in determining the initial guess, but are difficult to use for complete solution

We will now "guess" that the solution is $g(x) = \frac{1}{\sqrt{2\pi}\sigma_x} \exp\left\{-\frac{(x-\mu)^2}{2\sigma_x^2}\right\}$ P.6
 noting that this gaussian density has mean μ and variance σ_x^2 .
 Substituting this $g(x)$ into the Lemma above, we have that for any pdf $f(x)$

$$\begin{aligned} h(X) &= - \int_{-\infty}^{\infty} f(x) \log f(x) dx \leq - \int_{-\infty}^{\infty} f(x) \log g(x) dx \\ &= - \int_{-\infty}^{\infty} f(x) \left[-\log \sqrt{2\pi}\sigma_x - \frac{(x-\mu)^2}{2\sigma_x^2} \right] dx \\ &= \log \sqrt{2\pi}\sigma_x + \frac{\sigma_x^2}{2\sigma_x^2} = \frac{1}{2} \log 2\pi\sigma_x^2 + \frac{1}{2} \log e \\ &= \frac{1}{2} \log 2\pi e \sigma_x^2. \end{aligned}$$

$\therefore h(X) \leq \frac{1}{2} \log 2\pi e \sigma_x^2$ if X has mean μ and variance σ_x^2 , with equality iff $f(x) = g(x)$

\therefore The pdf $f(x)$ with mean μ and variance σ_x^2 that has maximum differential entropy $h(X)$ is

$$f(x) = \frac{1}{\sqrt{2\pi}\sigma_x} \exp\left\{-\frac{(x-\mu)^2}{2\sigma_x^2}\right\}$$

and the resulting ^{diff} entropy is

$$h_{\max}(X) = \frac{1}{2} \log 2\pi e \sigma_x^2.$$

h.b. One way to guess the answer is Gaussian is to use Lagrange multipliers.

$$\begin{aligned} \Phi(f(x)) &= - \int_{-\infty}^{\infty} f(x) \ln f(x) dx - \lambda_0 \left[\int_{-\infty}^{\infty} f(x) dx - 1 \right] - \lambda_1 \left[\int_{-\infty}^{\infty} x f(x) dx - \mu \right] \\ &\quad + \lambda_2 \left[\int_{-\infty}^{\infty} (x-\mu)^2 f(x) dx - \sigma_x^2 \right] \end{aligned}$$

$$\Rightarrow \phi(f(x)) = -f(x) \ln f(x) - \lambda_0 f(x) - \lambda_1 x f(x) - \lambda_2 f(x) (x-\mu)^2$$

$$\frac{\partial \phi}{\partial f(x)} = -\ln f(x) - 1 - \lambda_0 - \lambda_1 x - \lambda_2 (x-\mu)^2 = 0, \forall f(x)$$

$\Rightarrow f(x) = A \exp\{-\lambda_1 x - \lambda_2 (x-\mu)^2\}$. In principle, one could substitute $A, \lambda_1,$ and $\lambda_2,$ but this becomes very messy. But for $A = \frac{1}{\sqrt{2\pi}\sigma_x}, \lambda_1 = 0$ and $\lambda_2 = \frac{1}{2\sigma_x^2}$, this is the gaussian with mean μ and variance σ_x^2 , so this is a good (and the correct) guess.

7. Using Lagrange multipliers, we can set up the following objective function or Lagrangian for the $f(x)$ maximizing $h(X)$ subject to the constraints:

$$\Phi(f(x)) = - \int_{-\infty}^{\infty} f(x) \ln f(x) dx - \sum_{i=1}^n \lambda_i \left[\int_{-\infty}^{\infty} g_i(x) f(x) dx - \eta_i \right]$$

From this, we have the integrand objective function

$$\phi(f(x)) = -f(x) \ln f(x) - \sum_{i=1}^n \lambda_i g_i(x) f(x).$$

In order for $f(x)$ to be the maximizing pdf, from the Euler-Lagrange equation

$$\frac{\partial \phi(f(x))}{\partial f(x)} = 0, \quad \forall x \in \mathbb{R}.$$

So we have

$$\frac{\partial \phi(f(x))}{\partial f(x)} = -\ln f(x) + 1 - \sum_{i=1}^n \lambda_i g_i(x) = 0$$

$$\Rightarrow f(x) = \exp \left\{ 1 - \sum_{i=1}^n \lambda_i g_i(x) \right\} = A \exp \left\{ -\lambda_1 g_1(x) - \dots - \lambda_n g_n(x) \right\}$$

where A is a constant and $\lambda_1, \dots, \lambda_n$ (the Lagrange multipliers) are constants that can be found (in principle) by substitution back into the constraint equations

$$\int_{-\infty}^{\infty} f(x) g_i(x) dx = \eta_i, \quad i = 1, \dots, n.$$

Evaluating the differential entropy corresponding to this $f(x)$, we have

$$\begin{aligned} h(X) &= - \int_{-\infty}^{\infty} f(x) \ln f(x) dx = - \int_{-\infty}^{\infty} f(x) \ln \left[A \exp \left\{ -\lambda_1 g_1(x) - \dots - \lambda_n g_n(x) \right\} \right] dx \\ &= \int_{-\infty}^{\infty} f(x) \left[-\ln A + \lambda_1 g_1(x) + \dots + \lambda_n g_n(x) \right] dx \\ &= -\ln A + \lambda_1 \int_{-\infty}^{\infty} f(x) g_1(x) dx + \dots + \lambda_n \int_{-\infty}^{\infty} f(x) g_n(x) dx \\ &= -\ln A + \lambda_1 \eta_1 + \dots + \lambda_n \eta_n. \end{aligned}$$

8 (a)

$$\begin{aligned}
 -I(X;Y) &= - \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{f(x,y)}{f(x)f(y)} dx dy \\
 &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{f(x)f(y)}{f(x,y)} dx dy \stackrel{①}{\leq} \log \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \frac{f(x)f(y)}{f(x,y)} dx dy \\
 &= \log \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x)f(y) dx dy = \log \int_{-\infty}^{\infty} f(x) dx \int_{-\infty}^{\infty} f(y) dy = \log 1 = 0
 \end{aligned}$$

with equality iff $\frac{f(x,y)}{f(x)f(y)} = \text{const} = 1 \Leftrightarrow f(x,y) = f(x)f(y) \Leftrightarrow X \perp Y$

$\therefore I(X;Y) \geq 0$, with equality iff X and Y are statistically independent.

$$\begin{aligned}
 \text{(b)} \quad I(X;Y) &= h(X) - h(X|Y) = \int_{-\infty}^{\infty} f(x) \log \frac{1}{f(x)} dx - \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{1}{f(x|y)} dx dy \\
 &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{f(x|y)}{f(x)} dx dy = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{f(x,y)}{f(x)f(y)} dx dy \\
 &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{f(y|x)}{f(y)} dx dy \\
 &= \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{1}{f(y)} dx dy - \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f(x,y) \log \frac{1}{f(y|x)} dx dy \\
 &= h(Y) - h(Y|X) = I(Y;X)
 \end{aligned}$$

$$\therefore I(X;Y) = I(Y;X)$$

just as in the discrete case.

9. The cost-capacity function of the additive Gaussian noise channel with a mean power constraint is given by

$$C(\beta) = \sup_{f(x) \in \mathcal{F}(\beta)} \{I(X; Y)\}$$

where $\mathcal{F}(\beta)$ is the set of all pdfs such that $E[X^2] = \int_{\mathbb{R}} x^2 f(x) dx \leq \beta$.

We start by noting that

$$I(X; Y) = h(Y) - h(Y|X).$$

Now

$$h(Y|X) = E \left[\log \frac{1}{f(Y|X)} \right] = \iint_{\mathbb{R}^2} f_{XY}(x, y) \log \frac{1}{f(Y|X)} dy dx$$

$$= \int_{\mathbb{R}} f_X(x) \int_{\mathbb{R}} f(Y|X) \log \frac{1}{f(Y|X)} dy dx \dots (*)$$

Now $f(Y|X) = \frac{f_{XY}(x, y)}{f_X(x)}$

and

$$f_{XY}(x, y) = f_{XZ}(x(x, y), z(x, y)) \left| \frac{\partial(x, z)}{\partial(x, y)} \right|$$

where $x(x, y) = x$ and $z(x, y) = y - x$. Thus

$$\frac{\partial(x, z)}{\partial(x, y)} = \begin{vmatrix} \frac{\partial x}{\partial x} & \frac{\partial x}{\partial y} \\ \frac{\partial z}{\partial x} & \frac{\partial z}{\partial y} \end{vmatrix} = \begin{vmatrix} 1 & 0 \\ -1 & 1 \end{vmatrix} = 1$$

$$\therefore f_{XY}(x, y) = f_{XZ}(x, y-x) |1| = f_X(x) f_Z(y-x)$$

from which it follows that

$$f(Y|X) = \frac{f_X(x) f_Z(y-x)}{f_X(x)} = f_Z(y-x)$$

Now substituting into (*), we have

$$h(Y|X) = \int_{\mathbb{R}} f_X(x) \int_{\mathbb{R}} f_Z(y-x) \log \frac{1}{f_Z(y-x)} dy dx = \dots$$

$$= \int_{\mathbb{R}} f_x(x) \cdot h(z) dx = h(z)$$

because for any fixed x ,

$$\int_{\mathbb{R}} f_z(y-x) \log \frac{1}{f_z(y-x)} dy = \int_{\mathbb{R}} f_z(r) \log \frac{1}{f_z(r)} dr = h(z)$$

$$\therefore h(Y/X) = h(Z).$$

Let $\mathcal{F}(\beta) =$ set of all pdfs having mean 0 and variance less than or equal to β . Then

$$C(\beta) = \max_{f_x(x) \in \mathcal{F}(\beta)} \{I(X; Y)\} = \max_{f_x(x) \in \mathcal{F}(\beta)} \{h(Y) - h(Y/X)\}$$

$$= \max_{f_x(x) \in \mathcal{F}(\beta)} \{h(Y)\} - h(z)$$

$$= \max_{f_x(x) \in \mathcal{F}(\beta)} \{h(Y)\} - \frac{1}{2} \log 2\pi e \sigma_z^2 \quad \dots (*)$$

We now use the following Lemma to maximize $I(X; Y)$ and hence find $C(\beta)$:

Lemma: Let $f(x)$ and $g(x)$ be any two valid pdfs.

Then

$$\int_{\mathbb{R}} f(x) \log \frac{1}{g(x)} dx \geq \int_{\mathbb{R}} f(x) \log \frac{1}{f(x)} dx,$$

with equality iff $f(x) = g(x)$ (a.e.)

$$\text{Proof: } \int_{\mathbb{R}} f(x) \log \frac{1}{f(x)} dx - \int_{\mathbb{R}} f(x) \log \frac{1}{g(x)} dx = \int_{\mathbb{R}} f(x) \log \frac{g(x)}{f(x)} dx$$

$$\stackrel{①}{\leq} \log \left(\int_{\mathbb{R}} f(x) \frac{g(x)}{f(x)} dx \right) = \log \left(\int_{\mathbb{R}} g(x) dx \right) = \log 1 = 0$$

Furthermore, because $\log(\cdot)$ is strictly concave \uparrow , equality occurs iff $f(x) = g(x)$ a.e.

Claim

In order to find $C(\beta)$ as given in $(*)$, we must find the $f_x(x)$ having variance $\sigma_x^2 \leq \beta$ that maximizes $h(Y)$.

Claim: If X is a Gaussian with mean 0 and variance β , then Y is Gaussian with mean 0 and variance $\sigma_Y^2 = \beta + \sigma_Z^2$, and $I(X; Y)$ is maximized for all X with $E[X^2] \leq \beta$.

Proof: Because $X \perp Z$, Y will have variance $\sigma_Y^2 = \beta + \sigma_Z^2$ regardless of the distribution of X (under the constraint that $E[X^2] \leq \beta$).

Now consider any $f_Y(y)$ with mean 0 and variance $\sigma_Y^2 = \beta + \sigma_Z^2$, and let

$$g(y) = \frac{1}{\sigma_Y \sqrt{2\pi}} \exp\left(-\frac{y^2}{2\sigma_Y^2}\right). \text{ Then from the}$$

lemma above

$$\begin{aligned} h(Y) &= \int_{\mathbb{R}} f_Y(y) \log \frac{1}{f_Y(y)} dy \leq \int_{\mathbb{R}} f_Y(y) \log \frac{1}{g(y)} dy \\ &= \int_{\mathbb{R}} f_Y(y) \left[\log \sigma_Y \sqrt{2\pi} + \frac{y^2}{2\sigma_Y^2} \right] dy = \log \sigma_Y \sqrt{2\pi} + \frac{\sigma_Y^2}{2\sigma_Y^2} \\ &= \frac{1}{2} \log 2\pi e \sigma_Y^2 \quad \text{with equality iff} \\ & \quad f_Y(y) = \frac{1}{\sigma_Y \sqrt{2\pi}} \exp\left(-\frac{y^2}{2\sigma_Y^2}\right) \text{ (a.e.)} \end{aligned}$$

Thus it follows that

$$\begin{aligned} C(\beta) &= \frac{1}{2} \log 2\pi e \sigma_Y^2 - \frac{1}{2} \log 2\pi e \sigma_Z^2 \\ &= \frac{1}{2} \log 2\pi e (\sigma_Z^2 + \beta) - \frac{1}{2} \log 2\pi e \sigma_Z^2 \\ &= \frac{1}{2} \log \left(\frac{2\pi e (\sigma_Z^2 + \beta)}{2\pi e \sigma_Z^2} \right) = \frac{1}{2} \log \left(1 + \frac{\beta}{\sigma_Z^2} \right) \end{aligned}$$

and this is achieved by an X which is Gaussian with mean 0 and variance $\sigma_X^2 = \beta$.